
- The statement (preceding Theorem 3.2) that formulas in the original language cannot be tautologies unless they contain \(*\) has to be restricted to formulas containing at least one propositional variable, as the language is assumed to contain the truth constants 0, 1. (Thanks are due to Antonín Dvořák for pointing this out.)

- Some of the formal proofs needed for the completeness proof turned out to be circular; consequently, a few more axioms and derivation rules need be added to the axiomatic system given in the paper. (Hopefully, some of them are nonetheless redundant—we haven’t checked their independence yet.) The corrected axiomatic system has the following rules and axioms:

  \[
  \sigma T / \uparrow \sigma \varphi, \quad \text{for } T \vdash L \varphi \text{ and any substitution } \sigma \text{ of } S'\text{-formulae for atoms} \quad (L\uparrow) \\
  \varphi, \varphi \rightarrow \psi / \psi \quad (\text{mp}) \\
  \varphi, \varphi \equiv \psi / \psi \quad (\text{mp}_\equiv) \\
  0 / \varphi \quad (\text{efq}) \\
  \varphi / \uparrow \varphi \quad (!\text{-nec}) \\
  / !!(\varphi \equiv \psi) \quad (!\equiv) \\
  / \varphi \equiv \varphi \quad (\text{refl}_\equiv) \\
  (\varphi_1 \equiv \psi_1) \& \cdots \& (\varphi_n \equiv \psi_n) \rightarrow (c(\varphi_1, \ldots, \varphi_n) \equiv c(\psi_1, \ldots, \psi_n)), \quad (\text{egr}) \\
  \text{for each } n\text{-ary } c \in S' \\
  / !\varphi \& !\psi \rightarrow (((\varphi \equiv \psi) \leftrightarrow \uparrow \uparrow \Delta(\varphi \leftrightarrow \psi))) \quad (\equiv' \equiv) \\
  / !c(\varphi_1, \ldots, \varphi_n) \rightarrow !\varphi_i, \quad \text{for each } n\text{-ary } c \in S \text{ and } 1 \leq i \leq n \quad (c!) \\
  / !\varphi_1 \rightarrow \cdots \rightarrow (!\varphi_n \rightarrow !c(\varphi_1, \ldots, \varphi_n)) \cdots, \quad \text{for each } n\text{-ary } c \in S \quad (!c) \\
  / ((1 \rightarrow \varphi) \equiv 1) \rightarrow (\varphi \equiv 1) \quad (1\rightarrow) \\
  / !\varphi \vee (\varphi \equiv *) \quad (!/\vee) \\
  / (\varphi \equiv 1) \vee (\exists \varphi \equiv 0) \quad (\exists) \\
  / \uparrow * \quad (\uparrow *) \\
  / \neg !* \quad (?/*)